Self-similarity in the entropy graph for a family of piecewise linear maps.

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For fixed slope s > 1 and  $\gamma \in [0, 1]$ , take:

$$Q_\gamma(x) = egin{cases} x+1, & x\leq\gamma\ 1+s(1-x), & x>\gamma \end{cases}$$



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There is matching if there are  $\kappa^{\pm} \geq 1$  such that

$$egin{array}{rcl} Q^{\kappa^+}_\gamma(\gamma^+)&=&Q^{\kappa^-}_\gamma(\gamma^-),\ (Q^{\kappa^+}_\gamma)'(\gamma^+)&=&(Q^{\kappa^-}_\gamma)'(\gamma^-). \end{array}$$

The number

$$\Delta = \kappa^+ - \kappa^-$$

is the matching index.

Facts about matching for  $Q_{\gamma}$  (proved in our paper, not in this talk):

Matching holds for an open dense set of full Lebesgue measure.

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Matching holds for an open dense set of full Lebesgue measure.

 Matching occurs for every s-adic rational: γ = ps<sup>-q</sup> for p, q ∈ N. (In this case Q<sup>n</sup><sub>γ</sub>(γ<sup>+</sup>) = Q<sup>n</sup><sub>γ</sub>(γ<sup>-</sup>) = 1 is fixed for all n large enough.)

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- The non-matching set  $\mathcal{E}$  has Hausdorff dimension

$$\dim_H(\mathcal{E})=1,$$

but  $\dim_{H}(\mathcal{E} \setminus [0, \delta)) < 1$  for every  $\delta > 0$ .



Metric and topological entropy for s=2

Figure: Topological & metric entropies of  $Q_{\gamma}$  for s = 2 as functions of  $\gamma$ .

Theorem: Topological and metric entropy

$$h_{\mu}(Q_{\gamma}) ext{ and } h_{top}(Q_{\gamma}) ext{ are } \left\{ egin{array}{ll} ext{decreasing} & ext{if } \Delta < 0; \ ext{constant} & ext{if } \Delta = 0; \ ext{increasing} & ext{if } \Delta > 0, \end{array} 
ight.$$

as function of  $\gamma$  within matching intervals. (Recall  $\Delta = \kappa^+ - \kappa^-$ .)



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  - pseudo-centers
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- Tuning windows and explanation of the self-similarity of the entropy graphs.

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#### Why there is matching

Let  $g(x) := s(1-x) \mod 1$  and  $R : (0,1) \rightarrow (0,1)$  be the first return of  $Q_{\gamma}$  to [0,1).

Lemma:

$$R(x) = egin{cases} g(x) & ext{if } x \in (0,\gamma) \ g^2(x) & ext{if } x \in (\gamma,1) \end{cases}$$



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## Why there is matching



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Lemma: For fixed  $\gamma \in [0, 1]$ , the following conditions are equivalent:

(i) 
$$g^k(\gamma) < \gamma$$
 for some  $k \in \mathbb{N}$ ;

(ii) matching holds for  $\gamma$ .

In other words, the bifurcation set is

 $\mathcal{E} = \{ \gamma \in [0,1] : g^k(\gamma) \ge \gamma \ \forall k \in \mathbb{N} \}.$ 

Motivation: Find exact formulas for matching intervals J and their matching indices  $\Delta$  for slope s = 2 (also works for  $2 \le s \in \mathbb{N}$ ).

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Let  $\mathbb{Q}_{dyd}$  be the set of dyadic rationals in (0, 1].

Definition The pseudo-center of an interval  $J \subset (0, 1)$  is the (unique) dyadic rational  $\xi \in \mathbb{Q}_{dyd}$  with minimal denominator.

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#### Definition

- For binary string u, let  $\check{u}$  be the bitwise negation of u.
- For ξ ∈ Q<sub>dyd</sub> \ {1} and let w be the shortest even binary expansion of ξ and v be the shortest odd binary expansion of 1 − ξ.

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- Define the interval  $I_{\xi} := (\xi_L, \xi_R)$  containing  $\xi$  where,
- $\blacktriangleright \ \xi_L := . \overline{\check{v}v}, \quad \xi_R := . \overline{w}.$
- Also define the "degenerate" interval  $I_1 := (2/3, +\infty)$ .

In short: 
$$I_{\xi} := (\xi_L, \xi_R)$$
 with  $\xi_L := .\overline{v}v$ ,  $\xi_R := .\overline{w}$ .

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 with  $\xi_L := .\overline{vv}$ ,  $\xi_R := .\overline{w}$ .  
If  $\xi = 1/2$  then  $w = 10$ ,  $v = 1$  and  $\xi_L = .\overline{01}$ ,  $\xi_R = .\overline{10}$ .  
(01)  $w = u01 \Rightarrow \xi_L = .\overline{u001}\overline{u110}$ ;  
(11)  $w = u11 \Rightarrow \xi_L = .\overline{u101}\overline{u010}$ ;  
(010)  $w = u010 \Rightarrow \xi_L = .\overline{u00}\overline{u11}$ ;  
(110)  $w = u110 \Rightarrow \xi_L = .\overline{u10}\overline{u01}$ .

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Theorem:

- All matching intervals have the form *I<sub>ξ</sub>*, where *ξ* ∈ Q<sub>dyd</sub> are precisely the pseudo-centers of the components of [0, <sup>2</sup>/<sub>3</sub>] \ *E*.
- The matching index is

$$\Delta(\xi) = \frac{3}{2}(|w|_0 - |w|_1),$$

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Proposition: If  $\gamma \geq \frac{1}{6}$ , then  $|w_0| = |w|_1$ . In particular, all matching intervals in  $(\frac{1}{6}, \frac{2}{3})$  have matching index  $\Delta = 0$ 

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Remark: For general slope  $s \ge 2$ , the formulas are

$$\Delta(\xi) = \frac{s+1}{2} \sum_{a=0}^{s-1} (s-1-2a) |w|_a, \quad M = (\frac{s}{s+1} - \frac{1}{s}, \frac{s}{s+1})$$

Pseudo-center  $\xi = .w$  (even exp.) and  $1 - \xi = .v$  (odd exp.). The matching interval is  $I_{\xi} = [\xi_L, \xi_R]$  for  $\xi_L = .\overline{v}v$  and  $\xi_R = .\overline{w}$ .

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$$\xi_{\infty} \quad (\xi_2)_L = (\xi_3)_R \quad (\xi_1)_L = (\xi_2)_R \quad \xi_L = (\xi_1)_R \quad \xi \quad \xi_R$$

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Lemma: The pseudo-center of the next period doubling can be obtained from the previous using the substitution:

$$\chi: \begin{array}{ccc} w \mapsto \check{v}v & & \check{w} \mapsto v\check{v} \\ v \mapsto vw & & \check{v} \mapsto \check{v}\check{w} \end{array}.$$

Thus the limit  $\xi_{\infty}$  has *s*-adic expansion

 $\xi_{\infty} = .\check{v}\check{w}v\check{v}vw\check{v}\check{w}vw\check{v}v\check{v}\check{w}\dots$ 

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**Remark**: This proposition explains constant entropy on all matching intervals in  $M = (\frac{1}{6}, \frac{2}{3})$ . A no devil's staircase argument gives:

$$h_\mu(Q_\gamma) = \log(rac{1+\sqrt{5}}{2}) \quad ext{ and } h_{top}(Q_\gamma) = rac{2}{3}\log 2,$$

for all  $\gamma \in [\frac{1}{6}, \frac{2}{3}]$ .

## Pseudo-centers (tuning windows)

Pseudo-center  $\xi = .w$  (even expansion) and  $1 - \xi = .v$  (odd exp). The matching interval is  $I_{\xi} = [\xi_L, \xi_R] = [.\overline{v}v$ ,  $.\overline{w}]$ . The tuning interval is  $T_{\xi} = [\xi_T, \xi_R]$  for  $\xi_T = .v\overline{w}$ .



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Theorem: Let  $K(\xi_T) = \{x : g^k(x) \ge \xi_T \ \forall k\}$ . Then  $x \in K(\xi_T) \cap T_{\xi}$  if and only if  $x = .\sigma_1 \sigma_2 \sigma_3 \sigma_4 ...$ for  $\sigma_1 \in \{w, \check{v}\}, \sigma_j \in \{w, v, \check{w}, \check{v}\}$ describing a path in the diagram.  $\check{v}$  $\check{w}$ 

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# Shape of the entropy function



Question: Known:  $\gamma \mapsto h(\mu_{\gamma})$  is Hölder. Is  $\gamma \mapsto h_{top}(Q_{\gamma})$  Hölder?

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Conjecture: The neutral tuning windows are exactly the plateaus of (topological and metric) entropy.

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Corollary: The shape of the entire entropy function (i.e., pattern of increase/decrease) is repeated in every tuning window  $T_{\xi}$  with  $\Delta(\xi) > 0$ , and reversed in every tuning window  $T_{\xi}$  with  $\Delta(\xi) < 0$ .