# Primes as sums of Fibonacci numbers, II 

Lukas Spiegelhofer (MU Leoben)<br>Joint work with Michael Drmota and Clemens Müllner (TU Wien)

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We want to sketch the proof of the following theorem.
Proposition 1 (Drmota, Müllner, S. 2022+)
There exist constants $c>0$ and $C$ such that

$$
\sum_{p \leq x} \mathrm{e}(\vartheta \mathrm{z}(p)) \leq C(\log x)^{4} x^{1-c\|\vartheta\|^{2}}
$$

for all real $\vartheta$ and $x \geq 2$.

- $p$ is a prime,
- $\mathrm{e}(x)=\exp (2 \pi i x)$,
- $\mathrm{z}(n)$ is the minimal number of Fibonacci numbers needed to write $n$ as their sum,
- $\|\vartheta\|$ is the distance of $\vartheta$ to $\mathbb{Z}$.

The sum over primes can be rewritten, using summation by parts (e.g. Mauduit-Rivat 2010, Ann. of Math.):

$$
\sum_{p \leq N} \mathrm{e}(\vartheta \mathrm{z}(p)) \leq \frac{2}{\log N} \max _{t \leq N}\left|\sum_{n \leq t} \mathrm{e}(\vartheta \mathrm{z}(n)) \wedge(n)\right|+O(\sqrt{N})
$$

## Von Mangoldt function:

$$
\Lambda(n)= \begin{cases}\log p, & n=p^{k} \text { for some } k \geq 1 \text { and some prime } p \\ 0, & \text { otherwise }\end{cases}
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Vaughan's identity leads to the starting point of our method.
Lemma (cf. Davenport, Multiplicative number theory) Let $f: \mathbb{N} \rightarrow \mathbb{C}$ such that $|f(n)| \leq 1$ for all $n \geq 1$. For all $N, U, V \geq 2$ such that $U V \leq N$ we have

$$
\begin{aligned}
& \sum_{n \leq N} f(n) \Lambda(n) \ll U+(\log N) \sum_{t \leq U V} \max _{w}\left|\sum_{w \leq r \leq N / t} f(r t)\right| \\
& +\sqrt{N}(\log N)^{3} \max _{\substack{U \leq M \leq N / V \\
V \leq q \leq N / M}}\left(\sum_{V<p \leq N / M}\left|\sum_{\substack{M<m \leq 2 M \\
m \leq \min (N / p, N / q)}} f(m p) \overline{f(m q)}\right|\right)^{1 / 2}
\end{aligned}
$$

with an absolute implied constant.

Let us write

$$
\begin{aligned}
& S_{\mathrm{I}}(N, U, V):=\sum_{t \leq U V} \max _{w}\left|\sum_{w \leq r \leq N / t} f(r t)\right|, \\
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## Choose $U=N^{3 / 4}$ (and $V$ small).

- $S_{I}$, inner sum: the difference $t$ is usually large ( $\gg N^{3 / 4-\varepsilon}$ ) compared to the length of the sum $\left(\ll N^{1 / 4+\varepsilon}\right)$.
- $S_{\text {II }}$, inner sum: the differences $p$ and $q$ will be small $\left(\ll N^{1 / 4}\right)$ compared to the length of the sum $\left(\gg N^{3 / 4}\right)$.

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We establish a strong estimate for type-I sums -

## "the Zeckendorf sum-of-digits function has level of distribution 1".

Theorem (DMS 2022+)
Let $\varepsilon>0$. There exist $c_{1}=c_{1}(\varepsilon)>0$ and $C=C(\varepsilon)>0$ such that for all $\vartheta \in \mathbb{R}$ and all real $x \geq 2$,

$$
\sum_{1 \leq d \leq D} \max _{\substack{y, z \geq 0 \\ z-y \leq x}} \max _{0 \leq a<d}\left|\sum_{\substack{y \leq n<z \\ n \equiv a \bmod d}} \mathrm{e}(\vartheta \mathrm{z}(n))\right| \leq C(\log x)^{11 / 4} x^{1-c_{1}\|\vartheta\|^{2}}
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where $D=x^{1-\varepsilon}$.
This is a statement on z along very sparse finite arithmetic progressions

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A+\left(0, d^{\prime}, 2 d^{\prime}, 3 d^{\prime}, \ldots,(N-1) d^{\prime}\right)
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for average moduli $d$.

- The proof of this statement is based on the corresponding paper for the Thue-Morse sequence $\mathrm{t}(n)=(-1)^{\mathrm{s}_{2}(n)}$ (S., Compos. Math 2020):


## Sketch of proof of S2020.

- Apply the van der Corput inequality repeatedly in order to obtain higher order correlations.
- Estimate

nontrivially, where $\mathrm{s}_{2}^{(\rho)}(n)=s_{2}\left(n \bmod 2^{\rho}\right)$.
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## Eliminating digits

Consider a digital expansion along a sparse arithmetic progression: in each step, many digits change!

Each application of van der Corput's inequality "eliminates" digits with indices in a certain interval $[A, B)$.

It is easy to detect base- $q$ digits $\delta_{j}(n)$ with indices in $[A, B)$ : we have

$$
\left(\delta_{A}(n), \ldots, \delta_{B-1}(n)\right)=\left(\nu_{A}, \ldots, \nu_{B-1}\right) \quad \text { if and only if } \quad\left\{\frac{n}{q^{B}}\right\} \in J,
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where

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J=\left[\frac{m}{q^{B-A}}, \frac{m+1}{q^{B-A}}\right) \quad \text { and } \quad m=\sum_{A \leq j<B} \nu_{j} q^{j-A} .
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The indicator function $1 J$ can be approximated by trigonometric polynomials.

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## The Zeckendorf case: digit detection

Ostrowski expansion, $\varphi=\frac{\sqrt{5}+1}{2}$ : The Zeckendorf digits of $n$ with indices below $B$ are equal to prescribed values if and only if

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n \varphi
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is contained in a certain interval modulo 1 .


The blue and red intervals are separated $\leadsto$ integers having Zeckendorf expansion $\cdots * * 00 *$ cannot be detected by a single interval.

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## Wrap

Wrap this around the two-dimensional torus in such a way that adjacent digit combinations (w.r.t. the lexicographical ordering) lie "parallel to each other" (illustration in a moment): set

$$
p(n)=\left(\frac{n}{\varphi^{B}}, \frac{n}{\varphi^{B+1}}\right) .
$$

The closure of the set of points $p(n) \bmod 1 \times 1$ is a union of finitely many line segments, since

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Detecting the lowest Zeckendorf digit: $B=3$


The least significant digit is given by the Fibonacci word

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F=(0100101001001010010100100101001001 \cdots),
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## One significant digit: $\mathrm{B}=3$



Three significant digits
no longer separated!

$$
\underbrace{000,001}, 010,100,101
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The Zeckendorf digits of $n$ with indices in $[A, B)$ are equal to prescribed values $\omega_{A}, \ldots, \omega_{B-1}$ if and only if

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is contained in $Q+\mathbb{Z}^{2}$, where $Q$ is a certain parallelogram depending on the values $\omega_{j} . \quad \leadsto$ trigonometric approximation of $Q$ !

Together with the following lemma, we may eliminate Zeckendorf digits in our sum $\sum_{n} \mathrm{e}(\vartheta \mathrm{z}(n d+a))$.
Lemma (generalized vdC inequality)
Let I be a finite interval in $\mathbb{Z}$ containing $M$ integers and $x_{m} \in \mathbb{C}$ for $m \in I$. Assume that $K \subset \mathbb{N}$ is a finite nonempty set. Then


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- Apply this process repeatedly, until only few digits remain. - In analogy to the Thue-Morse case, it remained to estimate a Gowers (type) norm for the Zeckendorf sum-of-digits function.
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## Generalizations

For the (Ostrowski) $\alpha$-sum-of-digits function $\mathrm{s}_{\alpha}$, we expect that the method generalizes, at least for the case that $\alpha$ is an algebraic number. Algebraicity is probably needed for multi-dimensional detection.

More interestingly, we are interested in morphic sequences: take a fixed point of a substitution over a finite alphabet, and possibly rename the letters afterwards. Examples: the Thue-Morse sequence defined by
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## Linear recurrent number systems

The Tribonacci numeration system is based on the Tribonacci numbers

$$
\begin{aligned}
a_{n} & =a_{n-1}+a_{n-2}+a_{n-3}, \\
a_{0} & =a_{1}=0, a_{2}=1, \quad \text { that is, } \\
\left(a_{n}\right)_{n \geq 3} & =(1,2,4,7,13,24,44,81,149, \ldots) .
\end{aligned}
$$

Every natural number is the unique sum of pairwise different $a_{n}, n \geq 3$, where taking three consecutive numbers is forbidden.
The lowest two Tribonacci digits exhibit a close connection to the Tribonacci word


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a_{n} & =a_{n-1}+a_{n-2}+a_{n-3} \\
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\left(a_{n}\right)_{n \geq 3} & =(1,2,4,7,13,24,44,81,149, \ldots) .
\end{aligned}
$$

Every natural number is the unique sum of pairwise different $a_{n}, n \geq 3$, where taking three consecutive numbers is forbidden.
The lowest two Tribonacci digits exhibit a close connection to the Tribonacci word


## Linear recurrent number systems

The Tribonacci numeration system is based on the Tribonacci numbers

$$
\begin{aligned}
a_{n} & =a_{n-1}+a_{n-2}+a_{n-3}, \\
a_{0} & =a_{1}=0, a_{2}=1, \quad \text { that is, } \\
\left(a_{n}\right)_{n \geq 3} & =(1,2,4,7,13,24,44,81,149, \ldots) .
\end{aligned}
$$

Every natural number is the unique sum of pairwise different $a_{n}, n \geq 3$, where taking three consecutive numbers is forbidden.
The lowest two Tribonacci digits exhibit a close connection to the Tribonacci word

$$
\sigma:\left\{\begin{array}{lll}
0 & \mapsto & 01 \\
1 & \mapsto & 02 \\
2 & \mapsto & 0
\end{array}\right\} .
$$

## Tribonacci sum of digits

Meanwhile, the Tribonacci sum-of-digits function modulo 2 appears to be given by

$$
\sigma:\left\{\begin{array}{lll}
\mathrm{a} & \mapsto & \mathrm{ae} \\
\mathrm{~b} & \mapsto & \mathrm{af} \\
\mathrm{c} & \mapsto & \mathrm{a} \\
\mathrm{~d} & \mapsto & \mathrm{db} \\
\mathrm{e} & \mapsto & \mathrm{dc} \\
\mathrm{f} & \mapsto & \mathrm{~d}
\end{array}\right\}, \quad \pi:\left\{\begin{array}{lll}
\mathrm{a} & \mapsto & 0 \\
\mathrm{~b} & \mapsto & 0 \\
\mathrm{c} & \mapsto & 0 \\
\mathrm{~d} & \mapsto & 1 \\
\mathrm{e} & \mapsto & 1 \\
\mathrm{f} & \mapsto & 1
\end{array}\right\} .
$$

$$
\operatorname{tr}=0110100100101100101101101001011011010011010 \cdots
$$

Detecting a given letter in the Tribonacci word leads to the (classical) two-dimensional Rauzy fractal.


Figure: Jolivet, Loridant, Luo 2014

> For detecting Tribonacci digits with indices in $[A, B)$, we will have to consider three-dimensional cylinders with the Rauzy fractal as base!

Detecting a given letter in the Tribonacci word leads to the (classical) two-dimensional Rauzy fractal.


Figure: Jolivet, Loridant, Luo 2014
For detecting Tribonacci digits with indices in $[A, B)$, we will have to consider three-dimensional cylinders with the Rauzy fractal as base!

## Thank you!

