

The simplest 2-regular sequences

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Outline

1. k -regular sequences
2. The discrepancy of the van der Corput sequence
3. Odd binomial coefficients

Section 1

k -regular sequences

k-regular sequences

If there exists an integer r , sequences S_0, \dots, S_{r-1} , and $r \times r$ -matrices B_0, \dots, B_{k-1} such that

$$\begin{pmatrix} S_0(kn + a) \\ \vdots \\ S_{r-1}(kn + a) \end{pmatrix} = B_a \begin{pmatrix} S_0(n) \\ \vdots \\ S_{r-1}(n) \end{pmatrix}$$

for all $n \geq 0$ and $0 \leq a < k$, the sequence S_0 is called k -regular.
[Allouche–Shallit 1992]

First example. The *Stern sequence*, defined by

$$\begin{aligned} z(0) &= 0, & z(1) &= 1, \\ z(2n) &= z(n), \\ z(2n+1) &= z(n) + z(n+1) \end{aligned}$$

is 2-regular. Corresponding matrices are given by

$$B_0 = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}, \quad B_1 = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix},$$

and we have

$$\begin{aligned} \begin{pmatrix} z(2n) \\ z(2n+1) \end{pmatrix} &= B_0 \begin{pmatrix} z(n) \\ z(n+1) \end{pmatrix}, \\ \begin{pmatrix} z(2n+1) \\ z(2n+2) \end{pmatrix} &= B_1 \begin{pmatrix} z(n) \\ z(n+1) \end{pmatrix}. \end{aligned}$$

The values of the Stern sequence are distributed according to a log-normal distribution [Bettin, Drappeau, Spiegelhofer 2017]

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Second example

The binary sum-of-digits function satisfies

$$\begin{aligned}s(2n) &= s(n), \\ s(2n + 1) &= s(n) + 1.\end{aligned}$$

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Third example

The number A_n of odd entries in row n of Pascal's triangle satisfies

$$A_n = 2^{s(n)},$$

that is,

$$A_{2n} = A_n, \quad A_{2n+1} = 2A_n.$$

$$(A_{2n}) = (1) (A_n),$$

$$(A_{2n+1}) = (2) (A_n).$$

Fourth example

The *discrepancy of the van der Corput sequence* satisfies

$$d(1) = 1,$$

$$d(2n) = d(n),$$

$$d(2n+1) = \frac{d(n) + d(n+1) + 1}{2}.$$

$$\begin{pmatrix} d(2n) \\ d(2n+1) \\ 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & \frac{1}{2} \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} d(n) \\ d(n+1) \\ 1 \end{pmatrix},$$

$$\begin{pmatrix} d(2n+1) \\ d(2n+2) \\ 1 \end{pmatrix} = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} & \frac{1}{2} \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} d(n) \\ d(n+1) \\ 1 \end{pmatrix}.$$

Section 2

The discrepancy of the van der Corput sequence

It is known [Drmota, Larcher, Pillichshammer 2005] that $d(n)$ is normally distributed.

Theorem (DLP2005)

For every $y \in \mathbb{R}$ we have

$$\frac{1}{M} \# \left\{ N < M : d(n) \leq \frac{1}{4} \log_2 N + y \frac{1}{4\sqrt{3}} \sqrt{\log_2 N} \right\} = \Phi(y) + o(1),$$

where

$$\Phi(y) = \int_{-\infty}^y e^{-t^2/2} dt.$$

We want to obtain a finer distribution result using moments of $d(n)$, for example, involving a polynomially perturbed Gaussian

$$C \int_{-\infty}^y e^{-t^2/2} p(t) dt$$

together with a better error term.

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The expectation of $d(n)$

We have

$$\frac{1}{2^\lambda} \sum_{2^\lambda \leq n < 2^{\lambda+1}} d(n) = 1 + \frac{\lambda}{4}.$$

For $\lambda \geq 0$ and $2^\lambda \leq n < 2^{\lambda+1}$, we define

$$e_\lambda(n) = d(n) - 1 - \frac{\lambda}{4},$$

$$\Delta(n) = d(n) - d(n+1) = e_\lambda(n) - e_\lambda(n+1).$$

By construction,

$$\sum_{2^\lambda \leq n < 2^{\lambda+1}} e_\lambda(n) = \sum_{2^\lambda \leq n < 2^{\lambda+1}} \Delta(n) = 0.$$

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Some values

$$\begin{aligned}
 d(1) &= 1, & d(2) &= 1, & d(3) &= \frac{3}{2}, \\
 e_0(1) &= 0, & e_1(2) &= -\frac{1}{4}, & e_1(3) &= \frac{1}{4}, \\
 \Delta(1) &= 0, & \Delta(2) &= -\frac{1}{2}, & \Delta(3) &= \frac{1}{2},
 \end{aligned}$$

$$\begin{aligned}
 d(4) &= 1, & d(5) &= \frac{7}{4}, & d(6) &= \frac{3}{2}, & d(7) &= \frac{7}{4}, \\
 e_2(4) &= -\frac{1}{2}, & e_2(5) &= \frac{1}{4}, & e_2(6) &= 0, & e_2(7) &= \frac{1}{4}, \\
 \Delta(4) &= -\frac{3}{4}, & \Delta(5) &= \frac{1}{4}, & \Delta(6) &= -\frac{1}{4}, & \Delta(7) &= \frac{3}{4}.
 \end{aligned}$$

Mixed moments

For integers $k, \ell, \lambda \geq 0$ we define

$$m_{k,\ell,\lambda} = \frac{1}{2^\lambda} \sum_{2^\lambda \leq n < 2^{\lambda+1}} e_\lambda^k(n) \Delta^\ell(n).$$

After some rewriting we obtain

$$\begin{aligned} m_{k,\ell,\lambda+1} &= \frac{1}{2^{\lambda+1}} \sum_{2^\lambda \leq n < 2^{\lambda+1}} \left(e_\lambda(n) - \frac{1}{4} \right)^k \left(\frac{\Delta(n)}{2} - \frac{1}{2} \right)^\ell \\ &+ \frac{1}{2^{\lambda+1}} \sum_{2^\lambda \leq n < 2^{\lambda+1}} \left(e_\lambda(n) - \frac{\Delta(n)}{2} + \frac{1}{4} \right)^k \left(\frac{\Delta(n)}{2} + \frac{1}{2} \right)^\ell. \end{aligned}$$

The contribution of the moments $m_{k,r,\lambda}$, where $r \leq \ell$, is given by

$$\frac{1}{2^\ell} \sum_{\substack{0 \leq r \leq \ell \\ 2|r}} \binom{\ell}{r} m_{k,\ell-r,\lambda}.$$

We see that matrix products such as

$$\begin{pmatrix} \frac{1}{2^6} & 0 & \frac{1}{2^6} \binom{6}{2} & 0 & \frac{1}{2^6} \binom{6}{4} & 0 & \frac{1}{2^6} \binom{6}{6} \\ 0 & \frac{1}{2^5} & 0 & \frac{1}{2^5} \binom{5}{2} & 0 & \frac{1}{2^5} \binom{5}{4} & 0 \\ 0 & 0 & \frac{1}{2^4} & 0 & \frac{1}{2^4} \binom{4}{2} & 0 & \frac{1}{2^4} \binom{4}{4} \\ 0 & 0 & 0 & \frac{1}{2^3} & 0 & \frac{1}{2^3} \binom{3}{2} & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2^2} & 0 & \frac{1}{2^2} \binom{2}{2} \\ 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}^M \begin{pmatrix} m_{k,6,\lambda} \\ m_{k,5,\lambda} \\ m_{k,4,\lambda} \\ m_{k,3,\lambda} \\ m_{k,2,\lambda} \\ m_{k,1,\lambda} \\ m_{k,0,\lambda} \end{pmatrix}$$

will play a role!

Example.

$$A = \begin{pmatrix} \frac{1}{16} & 0 & \frac{3}{16} & 0 & \frac{1}{16} \\ 0 & \frac{1}{8} & 0 & \frac{3}{8} & 0 \\ 0 & 0 & \frac{1}{4} & 0 & \frac{1}{4} \\ 0 & 0 & 0 & \frac{1}{2} & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}, \quad \text{then}$$

$$A^k = \begin{pmatrix} \frac{1}{16^k} & 0 & \frac{2}{4^k} - \frac{2}{16^k} & 0 & \frac{16^k-1}{5 \cdot 16^k} - 2\frac{4^k-1}{3 \cdot 16^k} \\ 0 & \frac{1}{8^k} & 0 & \frac{1}{2^k} - \frac{1}{8^k} & 0 \\ 0 & 0 & \frac{1}{4^k} & 0 & \frac{16^k-4^k}{3 \cdot 16^k} \\ 0 & 0 & 0 & \frac{1}{2^k} & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}.$$

Idea: Formulas for powers of such matrices of size K will enable us to describe all moments of $e_\lambda(n)$ with precision $O(\lambda^{-K/2})$.



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Section 3

Odd binomial coefficients

Let us define

$$S(N) = \#\left\{ (n, k) : 0 \leq k \leq n < N : 2 \nmid \binom{n}{k} \right\} = \sum_{0 \leq n < N} 2^{s(n)}.$$

It is easy to see that

$$S(2N) = 3S(N), \quad S(2N + 1) = 2S(N) + S(N + 1).$$

Many authors have studied this quantity (Flajolet, Grabner, Harborth, Kirschenhofer, Larcher, Prodinger, Stolarsky, Tichy, ...).

In particular, $S(N)$ was compared to the function $f(x) = x^{\log 3 / \log 2}$, which also has the property $f(2x) = 3f(x)$, and therefore satisfies $f(2^N) = S(2^N)$.

There is a continuous function φ with period 1 such that

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The quantity

$$\begin{aligned} \min_{x \in [0,1]} \varphi(x) &= \liminf_{n \rightarrow \infty} S(N) 2^{-\log 3 / \log 2} \\ &= 0.8125565590160063876948821016495367124 \dots \end{aligned}$$

is called Stolarsky–Harborth constant.

Larcher (1992) showed that the position y of each minimum is given by a Sturmian word with slope $\vartheta = \log 3 / \log 2 - 1$: for some $d \in \mathbb{R}$, we have

$$y = \sum_{k \geq 1} 2^{-k} e_k,$$

where

$$e_k = \lfloor (k+1)\vartheta + d \rfloor - \lfloor k\vartheta + d \rfloor.$$

Admissible values of d , corresponding to minima, as well as their number, remain mysterious.

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Simplification of the problem

Let $\lambda \geq 0$ be an integer. On the interval $[2^\lambda, 2^{\lambda+1}]$ we compare S with the line segment connecting

$$(2^\lambda, 3^\lambda) \quad \text{and} \quad (2^{\lambda+1}, 3^{\lambda+1}).$$

Set

$$\Delta_\lambda(n) = S(N) - 3^\lambda - (n - 2^\lambda) \frac{2 \cdot 3^\lambda}{2^\lambda}.$$

Normalizing suitably, we obtain a continuous function

$$\Phi : [0, 1] \rightarrow \mathbb{R}$$

such that

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The Straight Stolarsky–Harborth constant


Surprisingly,

1. this does not seem to be handled in the literature,
2. the minimum of Φ seems to be attained at

$$x = \sum_{k \geq 1} 2^{-k} e_k = 0.67741147741147752252682561330016 \dots,$$

where

$$e_k = \lfloor (k+1)\vartheta + \mathbf{0} \rfloor - \lfloor k\vartheta + \mathbf{0} \rfloor$$

and $\vartheta = \log 3 / \log 2 - 1$ as before. 

The value of Φ at this position is

$$-0.4782436127025850978521771039475291 \dots$$

We might call this the Straight Stolarsky–Harborth constant.

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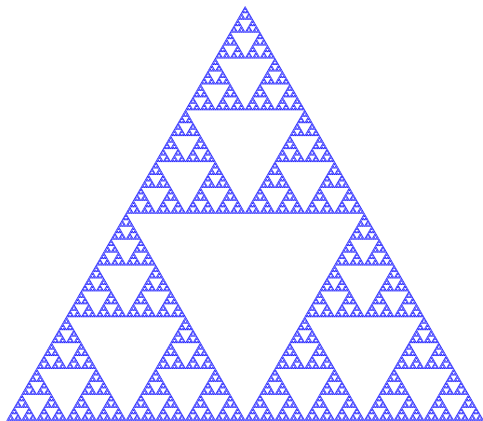
A remark

Note that the continuous function Φ satisfies, and is uniquely determined by,

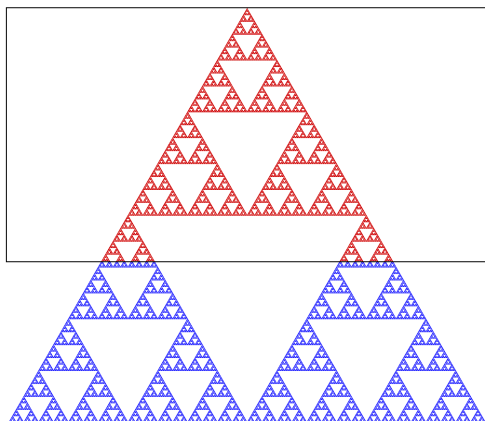
$$\begin{aligned}\Phi(0) &= \Phi(1) = 0, \\ \Phi\left(\frac{2k+1}{2^{\lambda+1}}\right) &= \frac{2}{3}\Phi\left(\frac{k}{2^\lambda}\right) + \frac{1}{3}\Phi\left(\frac{k+1}{2^\lambda}\right) - \frac{1}{3 \cdot 2^\lambda}\end{aligned}$$

for all integers $\lambda \geq 0$ and $2^\lambda \leq k < 2^{\lambda+1}$.

Remark. Consider the Sierpiński triangle T of height 1 and Hausdorff measure $H^{\log 3 / \log 2}(T) = C$, and the red portion $A(h)$ of height h .



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Prove that $H^{\log 3 / \log 2}(A(h)) = C\left(\frac{1}{2}\Phi(h) + h\right)$.

Higher divisibility in Pascal's triangle

Barat and Grabner (2001) also considered the number of binomial coefficients exactly divisible by a power of a prime.

Theorem (Barat–Grabner 2001, Theorem 5)

Let p be a prime, and $j \geq 0$ an integer. Define

$$S_j(N) = \# \left\{ (k, n) : 0 \leq k \leq n < N, p^j \parallel \binom{n}{k} \right\}.$$

There exist continuous periodic functions of period 1, $\psi_r^{(j)}$, for $0 \leq r \leq j$, such that

$$S_j(N) = N^\alpha \sum_{r=0}^j (\log_p N)^r \psi_r^{(j)}(\log_p N) + o(N^\varepsilon)$$

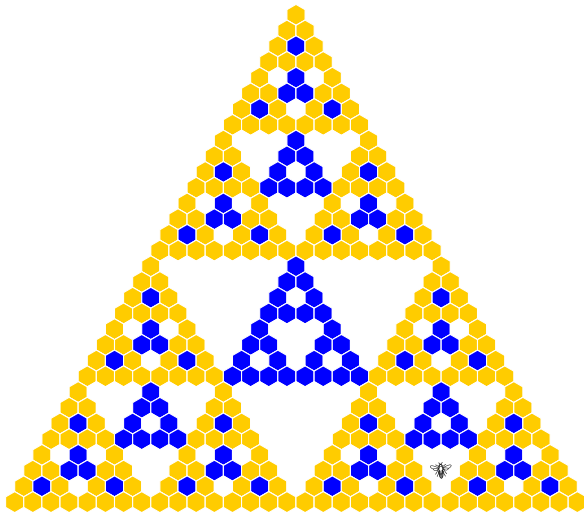
for $\alpha = \log_p \frac{p(p+1)}{2}$ and any $\varepsilon > 0$. Also, $\psi_j^{(j)} = \frac{1}{j!} \left(\frac{p-1}{p+1}\right)^{2j} \psi_0^{(0)}$.

More open questions

- ▶ (Grabner) Does there exist an exact representation of $S_j(N)$, where $j \geq 1$?
- ▶ What is the minimum of $\psi_r^{(j)}$ resp. of a *straightened* version?
- ▶ (Drmotá, Mauduit, Rivat) Can any bounds for the sum

$$\sum_{p \leq x} 2^{s(p)}$$

be proved?



THANK YOU!

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