# Thue-Morse along the sequence of cubes 

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## Section 1

## Subsequences of the Thue-Morse [tt: morrs] sequence

## The Thue-Morse sequence

We denote the Thue-Morse sequence on $\{0,1\}$ by $\mathbf{t}$. It is given by the binary sum-of-digits function $s$, reduced modulo 2 , or as the fixed point of the morphism $0 \mapsto 01,1 \mapsto 10$ that starts with 0 .


$$
\mathbf{t}=01101001100101101001011001101001 \cdots
$$

Also, denote the Thue-Morse sequence on $\{1,-1\}$ by $\mathbf{u}$.

## Thue-Morse $\rightleftharpoons$ Koch

Let $\mathrm{e}(x)=e^{2 \pi i x}$. The sequence $n \mapsto(-1)^{s(n)} \mathrm{e}(-n / 3)$ describes the direction of the $(n+1)$ th segment in the "unscaled Koch curve":


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## Arithmetic subsequences of $s$

Arithmetic subsequences of $\mathbf{t}$ are automatic.


Figure: An automaton for $\mathbf{t}(n)$


Figure: An automaton for $\mathbf{t}(3 n)$

The sum of digits along arithmetic progressions
For all integers $d \geq 0$ and rationals $x$ and $y$, the sequence $n \mapsto \mathrm{e}(s(d n) x+n y)$ is 2-automatic. Partial sums yield interesting pictures.


## The sum of digits along arithmetic progressions




Source: Wikipedia, "Rabbit"
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## The sum of digits along arithmetic progressions

The function $s_{q}$ along arithmetic progressions is uniformly distributed in residue classes modulo $m$ if $\operatorname{gcd}(q-1, m)=1$. We state the following special case.

Theorem (Gelfond 1968)
Let $d \geq 1$ and a be integers. There is an absolute $\lambda<1$ such that

$$
|\{1 \leq n \leq x: \mathbf{t}(n)=0, n \equiv a \bmod d\}|=\frac{x}{2 d}+\mathcal{O}\left(x^{\lambda}\right)
$$

## Very sparse arithmetic subsequences of $\mathbf{t}$

The Thue-Morse sequence has mean value around $1 / 2$ along most very short arithmetic progressions - $\mathbf{t}$ (and $\mathbf{u}$ ) has "level of distribution 1".

Theorem (S. 2020)
For all $\varepsilon>0$ we have

$$
\sum_{1 \leq d \leq D} \max _{1 \leq, z \geq 0} \max _{\substack{y, y \leq x}}\left|\sum_{\substack{y \leq n<z \\ n \equiv a \bmod d}} \mathbf{u}(n)\right| \leq C x^{1-\eta}
$$

for some $C$ and $\eta>0$ depending on $\varepsilon$, where $D=x^{1-\varepsilon}$.

In more relaxed language: let $R>0$. As $N \rightarrow \infty$, the following holds.

Most $d \asymp N^{R}$ have the property that

$$
\#\{0 \leq n<N: \mathbf{t}(n d+a)=0\}
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is close to $N / 2$ for all shifts $a$.

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## Problem $\stackrel{\text { 㘶 }}{\sim}$

Prove that for most $d \asymp N^{R}$,

$$
m \mapsto \#\{0 \leq n<N: s(n d+a)=m\}
$$


$2^{15}$ terms of $s\left(3^{21} n\right)$
closely follows a Gaussian for all shifts a.

## Polynomials of higher degree: Gelfond's third problem

Let $S=s_{q}$ be the sum-of-digits function in base $q \geq 2$.
Finalement, signalons comme problème à résoudre l'estimation du nombre des valeurs du polynôme $P(t)$ ne prenant que des valeurs entières sur l'ensemble [...] des entiers rationels, pour lesquelles on a $S[P(n)] \equiv \ell \bmod m$.
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That is, if $P$ is a polynomial such that $P(\mathbb{N}) \subseteq \mathbb{N}$, we are interested in

$$
A(q, P, m, \ell, x):=\#\left\{n<x: s_{q}(P(n)) \equiv \ell \bmod m\right\} .
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$$
\mathbf{t}=\left(s_{2}(n) \bmod 2\right)_{n \geq 0}
$$

$$
=\left(0^{\downarrow} 1^{\downarrow} 101^{\downarrow} 0011^{\downarrow} 0101101^{\downarrow} 001011001^{\downarrow} 10100110010{ }^{\downarrow} 1100110100101^{\downarrow} \cdots\right)
$$

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$$

$$
\begin{align*}
\mathbf{t} & =\left(s_{2}(n) \bmod 2\right)_{n \geq 0} \\
& =\left(\begin{array}{lll}
01 & 1 & 0
\end{array}\right. \tag{0}
\end{align*}
$$1

## Partial results

## We have

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$$
\begin{equation*}
\left|\#\left\{n<x: \mathbf{t}\left(n^{2}\right)=0\right\}-\frac{x}{2}\right| \leq C x^{1-c} . \tag{1}
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- For "sufficiently large bases" $q$ coprime to the leading coefficient of $P$, and $\operatorname{gcd}(q-1, m)=1$, the equivalence $A(q, P, m, \ell, x) \sim x / m$ has been proved (Drmota-Mauduit-Rivat 2011).


## The main result

Theorem (S. 2024+)
There exist real numbers $c>0$ and $C$ such that for all $x \geq 1$,

$$
\begin{equation*}
\left|\#\left\{n<x: \mathbf{t}\left(n^{3}\right)=0\right\}-\frac{x}{2}\right| \leq C x^{1-c} . \tag{2}
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Caveat. Currently, $c$ is not guaranteed to be greater than $2^{-500000}$.

## Section 2

## Sketch of the proof

## Carry Lemma (Mauduit-Rivat 2009, 2010)

- We are interested in the sum

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\sum_{n<2^{\nu}} \mathbf{u}\left((n+r)^{3}\right) \mathbf{u}\left(n^{3}\right)
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where $r$ is small compared to $2^{\nu}$.

- The arguments $(n+r)^{3}$ and $n^{3}$ usually have the same digits with indices above

$$
\lambda:=\nu(2+\varepsilon)
$$

- These digits can therefore be discarded.


## Carry lemma: a picture



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- In the correlation $\mathbf{u}\left((n+r)^{3}\right) \mathbf{u}\left(n^{3}\right)$, we may replace $\mathbf{u}$ by the restricted Thue-Morse sequence $\mathbf{u}^{[0, \lambda)}(n):=\mathbf{u}\left(n \bmod 2^{\lambda}\right)$.


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- The remaining window $[0, \lambda)$ of digits is about twice as long as the binary expansion of $n$.
- Therefore we cannot obtain uniform distribution of these digits, as $n$ runs.


## Too many significant digits

- A similar problem arises for sparse arithmetic progressions nd $+a$, where $n<2^{\nu}$ and $d \gg 2^{R \nu}$ : the binary digits of $(n+r) d+a$ and $n d+a$ usually differ up to index $\approx R \nu$.



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Wikipedia: "Salami", by Aka (CC BY-SA 2.5)

## Gowers norm

Iterated application (Cauchy-Schwarz, van der Corput) ${ }^{m}$ leads to higher order correlations of the Thue-Morse sequence, more precisely, a Gowers norm. This was estimated by Konieczny (2019) (and for general automatic sequences, by Byszewski-Konieczny-Müllner 2023): for some $c=c(m)>0$,

$$
\begin{aligned}
\frac{1}{2^{(m+1) \rho}} \sum_{n, r_{1}, \ldots, r_{m}<2^{\rho}} \prod_{\varepsilon_{1}, \ldots, \varepsilon_{m} \in\{0,1\}} \mathbf{u}^{[0, \rho)}(n+ & \left.\sum_{1 \leq i \leq m} \varepsilon_{i} r_{i}\right) \\
& =O(\exp (-c \rho)) .
\end{aligned}
$$

Iterated van der Corput could so far not be used for removing sufficiently many digits of polynomial values $P(n)$, if $\operatorname{deg} P>2$.

## A trivial decomposition ${ }^{1}$

- Choose $\rho<\nu$ in such a way that $3 \rho \geq \lambda$, and write

$$
n=2^{\rho} n_{1}+n_{0}, \quad \text { where } \quad\left\{\begin{array}{l}
0 \leq n_{1}<2^{\nu-\rho} \\
0 \leq n_{0}<2^{\rho}
\end{array}\right.
$$

- Expanding $n^{3} \bmod 2^{\lambda}$, we see that the cubic term in $n_{1}$ disappears.

[^0]




- On the critical interval $[2 \rho, \lambda)$ of length $\kappa:=\lambda-2 \rho$, the term $n_{1}^{2}$ is still relevant.

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- We introduce an additional sum $\sum_{0 \leq j<2^{\kappa}}$ that parametrizes the digit combinations in this interval.


## The critical interval of digits

We write

$$
\begin{aligned}
& S_{0}=\sum_{n<2^{\nu}} \mathbf{u}\left(n^{3}\right)=\sum_{n<2^{\nu}} \mathbf{u}^{[2 \rho, \lambda)}\left(n^{3}\right) \mathbf{u}^{\mathbb{N} \backslash[2 \rho, \lambda)}\left(n^{3}\right) \\
& \quad=\sum_{0 \leq j<2^{\kappa}} \mathbf{u}(j) \sum_{n<2^{\nu}} \mathbf{u}^{\mathbb{N} \backslash[2 \rho, \lambda)}\left(n^{3}\right) \llbracket \frac{n^{3}}{2^{\lambda}} \in\left[\frac{j}{2^{\kappa}}, \frac{j+1}{2^{\kappa}}\right)+\mathbb{Z} \| .
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\end{aligned}
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(1.) An additional sum of length $2^{\kappa}$ is introduced;
(2.) The rightmost factor detects whether $j$ corresponds to the digits of $n^{3}$ in the critical interval $\{\lambda-\kappa, \ldots, \lambda-1\}$.

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(1.) An additional sum of length $2^{\kappa}$ is introduced;
(2.) The rightmost factor detects whether $j$ corresponds to the digits of $n^{3}$ in the critical interval $\{\lambda-\kappa, \ldots, \lambda-1\}$.
(3.) After cutting away also the digits with indices $\geq \lambda$ (carry lemma),

A LINEAR SUM-OF-DIGITS PROBLEM IN $n_{1}$ REMAINS.

## Introducing trigonometric polynomials

- The "green" detecting term is approximated by a trigonometric polynomial $T$, evaluated at $\left(2^{\rho} n_{1}+n_{0}\right)^{3} / 2^{\lambda}$. The term $n_{1}^{3}$ does not appear in the argument of $T$.


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- Applying van der Corput's inequality another time, the argument of $T$ becomes linear in $n_{1}$, yielding a proper trigonometric polynomial in $n_{1}$ (cf. uniform distribution mod 1 of polynomial values).


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- Applying van der Corput's inequality another time, the argument of $T$ becomes linear in $n_{1}$, yielding a proper trigonometric polynomial in $n_{1}$ (cf. uniform distribution mod 1 of polynomial values).
- At this point all cubes and squares have been eliminated, at the cost of a much longer summation.


## Decoupling the trigonometric part

- The trigonometric polynomial in $n_{1}$ is decoupled from the sum over $n$, using suitable arithmetic subsequences and summation by parts. (Note that "everything is linear"!)


## Decoupling the trigonometric part

- The trigonometric polynomial in $n_{1}$ is decoupled from the sum over $n$, using suitable arithmetic subsequences and summation by parts. (Note that "everything is linear"!)
- The trigonometric component yields a geometric sum

$$
\varphi_{H}(x)=\sum_{0 \leq h<H} \mathrm{e}(h x) \ll \min \left(H,\|x\|^{-1}\right)
$$

where $\|x\|$ is the distance of $x$ to the nearest integer.
The average in $x$ of $\varphi_{H}(x)$ is only $\log H$ in size!

- Due to the small (logarithmic!) contribution of the critical interval, we only have to obtain a small gain in the sum-of-digits component.
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- Only arithmetic progressions play a role, which is is amenable to an iterated digit-elimination procedure [S2020]. This yields a gain $N^{-c}$ for some $c>0$, easily swallowing the logarithm.


## Essence of the proof

The additional sum introduced for digit detection in the critical interval only contributes a logarithm. A linear digital problem remains, which can be handled by iterated digit block elimination.

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Preprint，arXiv：2308．09498．


## THANK YOU!



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## van der Corput's inequality

## Lemma

Let I be a finite interval containing $N$ integers and let $a_{n}$ be a complex number for $n \in l$. For all integers $K \geq 1$ and $R \geq 1$ we have

$$
\left|\sum_{n \in I} a_{n}\right|^{2} \leq \frac{N+K(R-1)}{R} \sum_{|r|<R}\left(1-\frac{|r|}{R}\right) \sum_{\substack{n \in I \\ n+K r \in I}} a_{n+K r} \overline{a_{n}}
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$$

Instead of the original sum, we now have to estimate certain correlations (where $K R$ will be small compared to $N$ ).

## Higher degree polynomials

- Why not iterate the procedure of degree reduction?
- Note that

$$
\int_{0}^{1} \min \left(H,\|x\|^{-1}\right) \mathrm{d} x \asymp \log H
$$

while

$$
\int_{0}^{1}\left|\min \left(H,\|x\|^{-1}\right)\right|^{2} \mathrm{~d} x \asymp H
$$


[^0]:    ${ }^{1}$ Thanks to Michael Drmota, "maybe this can also be used for the cubes"

